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ESTIMATES FOR THE BEST CONSTANT  
IN A MARKOV  $L_2$ -INEQUALITY  
WITH THE ASSISTANCE OF COMPUTER ALGEBRA

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We prove two-sided estimates for the best (i.e., the smallest possible) constant  $c_n(\alpha)$  in the Markov inequality

$$\|p'_n\|_{w_\alpha} \leq c_n(\alpha) \|p_n\|_{w_\alpha}, \quad p_n \in \mathcal{P}_n.$$

Here,  $\mathcal{P}_n$  stands for the set of algebraic polynomials of degree  $\leq n$ ,  $w_\alpha(x) := x^\alpha e^{-x}$ ,  $\alpha > -1$ , is the Laguerre weight function, and  $\|\cdot\|_{w_\alpha}$  is the associated  $L_2$ -norm,

$$\|f\|_{w_\alpha} = \left( \int_0^\infty |f(x)|^2 w_\alpha(x) dx \right)^{1/2}.$$

Our approach is based on the fact that  $c_n^{-2}(\alpha)$  equals the smallest zero of a polynomial  $Q_n$ , orthogonal with respect to a measure supported on the positive axis and defined by an explicit three-term recurrence relation. We employ computer algebra to evaluate the seven lowest degree coefficients of  $Q_n$  and to obtain thereby bounds for  $c_n(\alpha)$ . This work is a continuation of a recent paper [5], where estimates for  $c_n(\alpha)$  were proven on the basis of the four lowest degree coefficients of  $Q_n$ .

**Keywords:** Markov type inequalities, Laguerre polynomials, three-term recurrence relation, Newton identities, computer algebra.

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## 1. INTRODUCTION AND STATEMENT OF THE RESULTS

Throughout this paper  $\mathcal{P}_n$  will stand for the set of algebraic polynomials of degree at most  $n$ , assumed, without loss of generality, with real coefficients. Let

$w_\alpha(x) := x^\alpha e^{-x}$ , where  $\alpha > -1$ , be the Laguerre weight function, and  $\|\cdot\|_{w_\alpha}$  be the associated  $L_2$ -norm,

$$\|f\|_{w_\alpha} = \left( \int_0^\infty |f(x)|^2 w_\alpha(x) dx \right)^{1/2}.$$

We study the best constant  $c_n(\alpha)$  in the Markov inequality in this norm

$$\|p'_n\|_{w_\alpha} \leq c_n(\alpha) \|p_n\|_{w_\alpha}, \quad p_n \in \mathcal{P}_n, \quad (1.1)$$

namely the constant

$$c_n(\alpha) := \sup_{p_n \in \mathcal{P}_n} \frac{\|p'_n\|_{w_\alpha}}{\|p_n\|_{w_\alpha}}.$$

Before formulating our results, let us give a brief account on the results known so far.

It is only the case  $\alpha = 0$  where the best Markov constant is known, namely, Turán [9] proved that

$$c_n(0) = \left( 2 \sin \frac{\pi}{4n+2} \right)^{-1}.$$

Dörfler [2] showed that  $c_n(\alpha) = \mathcal{O}(n)$  for every fixed  $\alpha > -1$  by proving the estimates

$$c_n^2(\alpha) \geq \frac{n^2}{(\alpha+1)(\alpha+3)} + \frac{(2\alpha^2+5\alpha+6)n}{3(\alpha+1)(\alpha+2)(\alpha+3)} + \frac{\alpha+6}{3(\alpha+2)(\alpha+3)}, \quad (1.2)$$

$$c_n^2(\alpha) \leq \frac{n(n+1)}{2(\alpha+1)}, \quad (1.3)$$

see [3] for a more accessible source. In the same paper, [3], Dörfler proved for the asymptotic constant

$$c(\alpha) := \lim_{n \rightarrow \infty} \frac{c_n(\alpha)}{n}, \quad (1.4)$$

that

$$c(\alpha) = \frac{1}{j_{(\alpha-1)/2,1}}, \quad (1.5)$$

where  $j_{\nu,1}$  is the first positive zero of the Bessel function  $J_\nu(z)$ .

Nikolov and Shadrin obtained in [5] the following result:

**Theorem A ([5, Theorem 1]).** *For all  $\alpha > -1$  and  $n \in \mathbb{N}$ ,  $n \geq 3$ , the best constant  $c_n(\alpha)$  in the Markov inequality (1.1) admits the estimates*

$$\frac{2(n + \frac{2\alpha}{3})(n - \frac{\alpha+1}{6})}{(\alpha+1)(\alpha+5)} < c_n^2(\alpha) < \frac{(n+1)(n + \frac{2(\alpha+1)}{5})}{(\alpha+1)[(\alpha+3)(\alpha+5)]^{1/3}}, \quad (1.6)$$

where for the left-hand inequality it is additionally assumed that  $n > (\alpha+1)/6$ .

Theorem A implies some inequalities for the asymptotic Markov constant  $c(\alpha)$  and, through (1.5), inequalities for  $j_{\nu,1}$ , the first positive zero of the Bessel function  $J_\nu$  (see [5, Corollaries 1,3]). It was also shown in [5, Theorem 2] that  $c(\alpha) = \mathcal{O}(\alpha^{-1})$ , which indicates that the upper estimate for  $c_n(\alpha)$  in Theorem A, though rather good for moderate  $\alpha$ , is not optimal.

In a recent paper [7] Nikolov and Shadrin proved an upper bound for  $c_n(\alpha)$  which is of the correct order with respect to both  $n$  and  $\alpha$  as they tend to infinity.

**Theorem B ([7, Theorem 1.1]).** *For all  $n \in \mathbb{N}$ ,  $n \geq 3$ , the best constant  $c_n(\alpha)$  in the Markov inequality (1.1) satisfies the inequality*

$$c_n^2(\alpha) \leq \frac{4n(n+2 + \frac{3(\alpha+1)}{4})}{\alpha^2 + 10\alpha + 8}, \quad \alpha \geq 2. \quad (1.7)$$

As a consequence of Theorem B and Dörfler's lower bound (1.2) for  $c_n(\alpha)$  Nikolov and Shadrin showed that

$$c_n^2(\alpha) \asymp \frac{n(n+\alpha+3)}{(\alpha+1)(\alpha+8)}, \quad n \geq 3, \alpha \geq 2.$$

**Corollary C ([7, Corollary 1.1]).** *For all  $\alpha \geq 2$  and  $n \geq 3$  the best constant  $c_n(\alpha)$  in the Markov inequality (1.1) satisfies*

$$\frac{2n(n+\alpha+3)}{3(\alpha+1)(\alpha+8)} \leq c_n^2(\alpha) \leq \frac{4n(n+\alpha+3)}{(\alpha+1)(\alpha+8)}. \quad (1.8)$$

In addition, Nikolov and Shadrin found the limit value of  $(\alpha+1)c_n^2(\alpha)$  as  $\alpha \rightarrow -1$ , and proved asymptotic inequalities for  $\alpha c_n^2(\alpha)$  as  $\alpha \rightarrow \infty$ .

**Corollary D ([7, Corollary 1.2]).** *The best constant  $c_n(\alpha)$  in the Markov inequality (1.1) satisfies:*

$$(i) \quad \lim_{\alpha \rightarrow -1} (\alpha+1)c_n^2(\alpha) = \frac{n(n+1)}{2};$$

$$(ii) \quad \frac{2n}{3} \leq \lim_{\alpha \rightarrow \infty} \alpha c_n^2(\alpha) \leq 3n.$$

A combination of Theorems A and B implies bounds for  $c(\alpha)$  defined in (1.4):

**Corollary E ([7, Corollary 1.3]).** *The asymptotic Markov constant  $c(\alpha)$  satisfies*

$$\frac{2}{(\alpha+1)(\alpha+5)} < c^2(\alpha) < \begin{cases} \frac{1}{(\alpha+1)\sqrt[3]{(\alpha+3)(\alpha+5)}}, & -1 < \alpha \leq \alpha^*, \\ \frac{1}{\alpha^2 + 10\alpha + 8}, & \alpha > \alpha^*, \end{cases}$$

where  $\alpha^* \approx 43.4$ .

The ratio of the upper and the lower bound for  $c(\alpha)$  in Corollary E is less than  $\sqrt{2}$  for all  $\alpha > -1$ .

In this paper we investigate the best Markov constant  $c_n(\alpha)$  following the approach from [5]. It is known (see Proposition 1 below) that  $c_n^{-2}(\alpha)$  is equal to the smallest zero of a polynomial  $Q_n$ , which is orthogonal with respect to a measure supported on  $\mathbb{R}_+$ . Since  $\{Q_n\}_{n \in \mathbb{N}}$  are defined by an explicit three-term recurrence relation, one can evaluate (at least theoretically) as many coefficients of  $Q_n$  as necessary. With the assistance of Wolfram's *Mathematica* we find the seven lowest degree coefficients of the polynomial  $Q_n$ , and thereby the six highest degree coefficients of  $R_n$ , the monic polynomial reciprocal to  $Q_n$ . Then we apply a simple technique for estimating the largest zero  $x_n$  of  $R_n$  on the basis of its  $k$  highest degree coefficients,  $3 \leq k \leq 6$ , thus obtaining lower and upper bounds for  $c_n^2(\alpha)$ . Our main result in this paper is:

**Theorem 1.** For  $3 \leq k \leq 6$  and for all  $n \geq k$ , the best constant  $c_n(\alpha)$  in the Markov inequality (1.1) admits the estimates

$$\underline{c}_{n,k}(\alpha) \leq c_n(\alpha) \leq \bar{c}_{n,k}(\alpha), \quad \alpha > -1, \quad (1.9)$$

where

$$\underline{c}_{n,3}^2(\alpha) = \frac{2n(n + \frac{3(\alpha+1)}{8})}{(\alpha+1)(\alpha+5)}, \quad (1.10)$$

$$\bar{c}_{n,3}^2(\alpha) = \frac{(n+1)(n + \frac{2(\alpha+1)}{5})}{(\alpha+1)[(\alpha+3)(\alpha+5)]^{1/3}}, \quad (1.11)$$

$$\underline{c}_{n,4}^2(\alpha) = \frac{(5\alpha+17)n(n + \frac{8(\alpha+1)}{25})}{2(\alpha+1)(\alpha+3)(\alpha+7)}, \quad (1.12)$$

$$\bar{c}_{n,4}^2(\alpha) = \frac{(5\alpha+17)^{1/4}(n+1)(n + \frac{3(\alpha+1)}{7})}{(\alpha+1)(\alpha+3)^{1/2}[2(\alpha+5)(\alpha+7)]^{1/4}}, \quad (1.13)$$

$$\underline{c}_{n,5}^2(\alpha) = \frac{2(7\alpha+31)n(n + \frac{25(\alpha+1)}{84})}{(\alpha+1)(\alpha+9)(5\alpha+17)}, \quad (1.14)$$

$$\bar{c}_{n,5}^2(\alpha) = \frac{(7\alpha+31)^{1/5}(n+1)(n + \frac{4(\alpha+1)}{9})}{(\alpha+1)(\alpha+3)^{2/5}[(\alpha+5)(\alpha+7)(\alpha+9)]^{1/5}}, \quad (1.15)$$

$$\underline{c}_{n,6}^2(\alpha) = \frac{(21\alpha^3 + 299\alpha^2 + 1391\alpha + 2073)n(n + \frac{2(\alpha+1)}{7})}{(\alpha+1)(\alpha+3)(\alpha+5)(\alpha+11)(7\alpha+31)}, \quad (1.16)$$

$$\bar{c}_{n,6}^2(\alpha) = \frac{(21\alpha^3 + 299\alpha^2 + 1391\alpha + 2073)^{1/6}(n+1)(n + \frac{5(\alpha+1)}{11})}{(\alpha+1)(\alpha+3)^{1/2}(\alpha+5)^{1/3}[(\alpha+7)(\alpha+9)(\alpha+11)]^{1/6}}. \quad (1.17)$$

**Remark 1.** For  $3 \leq k \leq 6$ , the pair  $(\underline{c}_{n,k}(\alpha), \bar{c}_{n,k}(\alpha))$  of bounds for  $c_n(\alpha)$  is deduced with the use of the  $k$  highest degree coefficients of the polynomial  $R_n$  (and (1.11) is also proved in [5]). Generally, the bounds for  $c_n(\alpha)$  obtained with larger  $k$  are better, though some exceptions are observed for small  $n$  and  $\alpha$ .

Clearly, inequalities (1.9) imply bounds for the asymptotic Markov constant  $c(\alpha)$ . Here, it is not difficult to prove that the larger  $k$ , the better the implied lower and upper bounds for  $c(\alpha)$ , hence the best bounds for  $c(\alpha)$  are obtained from (1.9) with  $k = 6$ .

Thus, Theorem 1 yields an improvement of the estimates for the asymptotic Markov constant  $c(\alpha)$  in Corollary E.

**Corollary 1.** *The asymptotic Markov constant  $c(\alpha) = \lim_{n \rightarrow \infty} n^{-1}c_n(\alpha)$  satisfies the inequalities*

$$\underline{c}(\alpha) < c(\alpha) < \bar{c}(\alpha),$$

where

$$\underline{c}^2(\alpha) := \frac{21\alpha^3 + 299\alpha^2 + 1391\alpha + 2073}{(\alpha + 1)(\alpha + 3)(\alpha + 5)(\alpha + 11)(7\alpha + 31)}$$

and

$$\bar{c}^2(\alpha) := \begin{cases} \frac{(21\alpha^3 + 299\alpha^2 + 1391\alpha + 2073)^{1/6}}{(\alpha + 1)(\alpha + 3)^{1/2}(\alpha + 5)^{1/3}[(\alpha + 7)(\alpha + 9)(\alpha + 11)]^{1/6}}, & -1 < \alpha \leq \alpha^*, \\ \frac{4}{\alpha^2 + 10\alpha + 8}, & \alpha > \alpha^*, \end{cases}$$

with  $\alpha^* \approx 172$ .

It is worth noticing that the ratio of the upper and the lower bound for  $c(\alpha)$  in Corollary 1 does not exceed  $\frac{2\sqrt{3}}{3} \approx 1.1547$  for all  $\alpha > -1$ .

Theorem 1, in particular inequality (1.16), implies an improvement of the lower bound in Corollary D(ii).

**Corollary 2.** *The best constant  $c_n(\alpha)$  in the Markov inequality (1.1) satisfies:*

$$\frac{6n}{7} \leq \lim_{\alpha \rightarrow \infty} \alpha c_n^2(\alpha) \leq 3n.$$

The rest of the paper is organized as follows. Section 2 contains some preliminaries. In Section 2.1 we characterize the squared best Markov constant as the largest zero of an  $n$ -th degree monic polynomial  $R_n$  with positive roots, and propose a recursive procedure for the evaluation of its coefficients (Proposition 2). Two-sided estimates for the largest zero of polynomials with only positive roots in terms of few of their coefficients are proposed in Sect. 2.2 (Proposition 2.3). The assisted by Wolfram's *Mathematica* proof of our results is given in Section 3.

In Section 4 we give some final remarks and conclusions, and formulate two conjectures concerning the asymptotic behavior of the best Markov constant and the coefficients of the characteristic polynomial  $R_n$ .

## 2. PRELIMINARIES

### 2.1. AN ORTHOGONAL POLYNOMIAL RELATED TO $c_n(\alpha)$

It is well-known that the squared best constant in a Markov-type inequality in  $L_2$ -norm is equal to the largest eigenvalue of a related positive definite  $n \times n$  matrix  $\mathbf{A}_n$ , thus the problem of finding the best Markov constant is equivalent to evaluating the largest eigenvalue of  $\mathbf{A}_n$ . Perhaps, a less known fact is that for a wide class of  $L_2$ -norms, the inverse matrix  $\mathbf{A}_n^{-1}$  is tri-diagonal, see [1, Sect. 2]. In the particular case of the  $L_2$ -norm induced by the Laguerre weight function  $w_\alpha$  this connection is given by the following proposition:

**Proposition 1** ([3, p. 85]). *The quantity  $c_n^{-2}(\alpha)$  is equal to the smallest zero of the polynomial  $Q_n(x) = Q_n(x, \alpha)$ , which is defined recursively by*

$$Q_{n+1}(x) = (x - d_n)Q_n(x) - \lambda_n^2 Q_{n-1}(x), \quad n \geq 0;$$

$$Q_{-1}(x) := 0, \quad Q_0(x) := 1;$$

$$d_0 := 1 + \alpha, \quad d_n := 2 + \frac{\alpha}{n+1}, \quad n \geq 1;$$

$$\lambda_0 > 0 \text{ arbitrary}, \quad \lambda_n^2 := 1 + \frac{\alpha}{n}, \quad n \geq 1.$$

By Favard's theorem, for any  $\alpha > -1$ ,  $\{Q_n(x, \alpha)\}_{n=0}^\infty$  form a system of monic orthogonal polynomials. Since  $Q_n$  is the characteristic polynomial of the inverse of a positive definite matrix (which is also positive definite), it follows that all the zeros of  $Q_n$  are positive (and distinct). Consequently,  $\{Q_n\}_{n=0}^\infty$  are orthogonal with respect to a measure supported on  $\mathbb{R}_+$ .

By Proposition 1, we have

$$Q_{n+1}(x) = \left(x - 2 - \frac{\alpha}{n+1}\right)Q_n(x) - \left(1 + \frac{\alpha}{n}\right)Q_{n-1}(x), \quad n \geq 1, \quad (2.1)$$

$$Q_0(x) = 1, \quad Q_1(x) = x - \alpha - 1. \quad (2.2)$$

If we write  $Q_n$  in the form

$$Q_n(x) = x^n - a_{n-1,n}x^{n-1} + a_{n-2,n}x^{n-2} - \dots + (-1)^n a_{0,n},$$

then

$$a_{0,n} = \binom{n+\alpha}{n}, \quad n \in \mathbb{N}_0, \quad (2.3)$$

with the convention that the right-hand side is equal to 1 for  $n = 0$ . The proof is by induction with respect to  $n$ . For  $n = 0, 1$ , (2.3) follows from (2.2). Assuming (2.3) is true for all  $m \leq n$ , we verify it for  $m = n + 1$  by putting  $x = 0$  in (2.1) and using the induction hypothesis:

$$\begin{aligned} (-1)^{n+1}a_{0,n+1} &= \left(2 + \frac{\alpha}{n+1}\right)(-1)^{n+1}\binom{n+\alpha}{n} + \left(1 + \frac{\alpha}{n}\right)(-1)^n\binom{n-1+\alpha}{n-1} \\ &= (-1)^{n+1}\binom{n+1+\alpha}{n}. \end{aligned}$$

Now, instead of  $\{Q_n\}_{n=0}^\infty$ , we consider the sequence of orthogonal polynomials  $\{\tilde{Q}_n\}_{n=0}^\infty$  normalized so that  $\tilde{Q}_n(0) = 1$ ,  $n \in \mathbb{N}_0$ , i.e.,

$$Q_n(x) = (-1)^n \binom{n+\alpha}{n} \tilde{Q}_n(x), \quad n \in \mathbb{N}_0.$$

It follows from (2.1) and (2.2) that  $\{\tilde{Q}_n\}_{n \in \mathbb{N}_0}$  are determined by

$$\left(1 + \frac{\alpha}{n+1}\right)\tilde{Q}_{n+1}(x) = \left(2 + \frac{\alpha}{n+1} - x\right)\tilde{Q}_n(x) - \tilde{Q}_{n-1}(x), \quad n \geq 1, \quad (2.4)$$

$$\tilde{Q}_0(x) = 1, \quad \tilde{Q}_1(x) = 1 - \frac{x}{\alpha+1}. \quad (2.5)$$

Writing  $\tilde{Q}_n$  in the form

$$\tilde{Q}_n(x) = 1 - A_{1,n}x + A_{2,n}x^2 - \cdots + (-1)^n A_{n,n}x^n$$

and rewriting (2.4) as

$$\tilde{Q}_{n+1}(x) - \tilde{Q}_n(x) = \frac{n+1}{n+\alpha+1}(\tilde{Q}_n(x) - \tilde{Q}_{n-1}(x)) + \frac{n+1}{n+\alpha+1}x\tilde{Q}_n(x), \quad n \in \mathbb{N},$$

we deduce the following recurrence relation for the evaluation of the coefficients  $\{A_{i,m}\}$ :

$$A_{i,n+1} - A_{i,n} = \frac{n+1}{n+\alpha+1}(A_{i,n} - A_{i,n-1}) + \frac{n+1}{n+\alpha+1}A_{i-1,n}, \quad n \geq k \geq 1, \quad (2.6)$$

with  $A_{0,n} = 1$  and  $A_{1,1} = \frac{1}{\alpha+1}$ .

Since, by Proposition 1,  $c_n^{-2}(\alpha)$  is equal to the smallest zero of  $\tilde{Q}_n$ , it follows that  $c_n^2(\alpha)$  equals the largest zero of the reciprocal polynomial of  $\tilde{Q}_n$ ,

$$R_n(x) = x^n \tilde{Q}_n(1/x). \quad (2.7)$$

The above observations allow us to reformulate Proposition 1 in the following equivalent form:

**Proposition 2.** *The squared best Markov constant  $c_n^2(\alpha)$  is equal to the largest zero of the polynomial*

$$R_n(x) = x^n - A_{1,n}x^{n-1} + A_{2,n}x^{n-2} - \dots + (-1)^n A_{n,n}. \quad (2.8)$$

The coefficients of  $R_n$  are evaluated recursively by the following procedure:

- $A_{1,1} = \frac{1}{\alpha+1}$ ;
- Set  $A_{0,m} = 1$ ,  $m = 0, \dots, n$ ;
- For  $i = 1$  to  $n$ :

1. Find the sequence  $\{D_{i,m}\}_{m=i-1}^n$  as solution of the recurrence equation

$$D_{i,m+1} = \frac{m+1}{m+\alpha+1} D_{i,m} + \frac{m+1}{m+\alpha+1} A_{i-1,m} \quad (2.9)$$

with the initial condition  $D_{i,i-1} = 0$ ;

2. Evaluate

$$A_{i,n} = \sum_{m=i}^n D_{i,m}. \quad (2.10)$$

## 2.2. POLYNOMIALS WITH POSITIVE ROOTS: BOUNDS FOR THE LARGEST ZERO

Let  $P$  be a monic polynomial of degree  $n$  with zeros  $\{x_i\}_{i=1}^n$ ,

$$P(x) = \prod_{i=1}^n (x - x_i) = x^n - b_1 x^{n-1} + b_2 x^{n-2} - \dots + (-1)^n b_n.$$

The coefficients  $b_r = b_r(P)$ ,  $r = 1, \dots, n$ , are given by the elementary symmetric functions of  $\{x_i\}_{i=1}^n$ ,

$$b_r = s_r = s_r(P) = \sum_{1 \leq i_1 < i_2 < \dots < i_r \leq n} x_{i_1} x_{i_2} \dots x_{i_r}, \quad r = 1, \dots, n.$$

It is well known that the elementary symmetric functions  $\{s_r\}$  and the Newton functions (sums of powers of  $x_i$ )

$$p_r = p_r(P) = \sum_{i=1}^n x_i^r, \quad r = 1, 2, 3, \dots,$$

are connected by the Newton identities:

$$p_r + \sum_{i=1}^{r-1} (-1)^i p_{r-i} s_i + (-1)^r r s_r = 0, \quad \text{if } 1 \leq r \leq n, \quad (2.11)$$

$$p_r + \sum_{i=1}^n (-1)^i p_{r-i} s_i = 0, \quad \text{if } r > n. \quad (2.12)$$



For a proof, see e.g. [10] or [4].

Our interest in the Newton functions is motivated by the fact that they provide tight bounds for the largest zero of a polynomial whose roots are all positive. For any such polynomial  $P$ , we set

$$\ell_k(P) := \frac{p_k(P)}{p_{k-1}(P)}, \quad u_k(P) := [p_k(P)]^{1/k}, \quad k \in \mathbb{N},$$

with the convention that  $p_0(P) := \deg(P)$ .

**Proposition 3.** *Let  $P(x) = x^n - b_1 x^{n-1} + b_2 x^{n-2} - \dots + (-1)^{n-1} b_{n-1} x + (-1)^n b_n$  be a polynomial with positive zeros  $x_1 \leq x_2 \leq \dots \leq x_n$ .*

*Then the largest zero  $x_n$  of  $P$  satisfies the inequalities*

$$\ell_k(P) \leq x_n < u_k(P), \quad k \in \mathbb{N}. \quad (2.13)$$

*Moreover, the sequence  $\{\ell_k(P)\}_{k=1}^\infty$  is monotonically increasing while the sequence  $\{u_k(P)\}_{k=1}^\infty$  is monotonically decreasing, and*

$$\lim_{k \rightarrow \infty} \ell_k(P) = \lim_{k \rightarrow \infty} u_k(P) = x_n. \quad (2.14)$$

*Proof.* For  $i = 1, \dots, n-1$ , we set  $a_i := \frac{x_i}{x_n}$ , then  $0 < a_i \leq 1$ . Now both inequalities (2.13) and the limit relations (2.14) readily follow from the representations

$$\ell_k(P) = \frac{a_1^k + \dots + a_{n-1}^k + 1}{a_1^{k-1} + \dots + a_{n-1}^{k-1} + 1} x_n, \quad u_k(P) = (a_1^k + \dots + a_{n-1}^k + 1)^{1/k} x_n.$$

The monotonicity of the sequence  $\{\ell_k(P)\}_{k=1}^\infty$  follows easily from Cauchy-Bouniakowsky's inequality. Indeed, we have

$$\left( \sum_{i=1}^n x_i^k \right)^2 = \left( \sum_{i=1}^n x_i^{\frac{k-1}{2}} x_i^{\frac{k+1}{2}} \right)^2 \leq \left( \sum_{i=1}^n x_i^{k-1} \right) \left( \sum_{i=1}^n x_i^{k+1} \right),$$

whence  $p_k^2(P) \leq p_{k-1}(P) p_{k+1}(P)$ , and consequently

$$\ell_k(P) = \frac{p_k(P)}{p_{k-1}(P)} \leq \frac{p_{k+1}(P)}{p_k(P)} = \ell_{k+1}(P).$$

To prove monotonicity of the sequence  $\{u_k(P)\}_{k=1}^\infty$ , we recall that  $0 < a_i \leq 1$  and therefore  $a_i^{k+1} \leq a_i^k$ . We have

$$(a_1^{k+1} + \dots + a_{n-1}^{k+1} + 1)^{1/(k+1)} < (a_1^{k+1} + \dots + a_{n-1}^{k+1} + 1)^{1/k} \leq (a_1^k + \dots + a_{n-1}^k + 1)^{1/k},$$

which yields

$$u_{k+1}(P) < u_k(P).$$

□

### 3. COMPUTER ALGEBRA ASSISTED PROOF OF THE RESULTS

Here we give the algorithms, the source code and the results of the computer algebra assisted proof of estimates (1.10)-(1.17) in Theorem 1. While the case  $k = 3$  and to a certain extent  $k = 4$  could be studied by hand, it seems impossible to provide similar calculations for larger  $k$ . We implement the idea from [5] for estimating  $c_n(\alpha)$  using  $k = 3$  highest degree coefficients of the polynomial  $R_n(x)$  and with the assistance of Wolfram's *Mathematica* v. 10 software we investigate the cases  $k = 4, 5, 6$ , as well. Software based on the algorithms described below failed with calculations for  $k > 6$ .

Henceforth, we write the polynomial  $R_n$  from (2.7) and (2.8) in the form

$$R_n(x) = x^n - b_1x^{n-1} + b_2x^{n-2} + \dots + (-1)^n b_n.$$

#### 3.1. LOWER BOUNDS FOR $c_n(\alpha)$

We apply Proposition 3 to estimate the largest zero  $x_n = c_n^2(\alpha)$  of the polynomial  $R_n(x)$  from below,

$$x_n \geq \ell_k(R_n) = \frac{p_k(R_n)}{p_{k-1}(R_n)}, \quad k = 3, 4, 5, 6,$$

and then with the help of computer algebra obtain a further estimation of the form

$$\ell_k(R_n) \geq cn(n + \sigma(\alpha + 1)),$$

with the optimal (i.e., the largest possible) constants  $c = c(k)$  and  $\sigma = \sigma(k)$ .

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**Algorithm 1** Estimating  $c_n(\alpha)$  from below

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- Input:*  $k \in \{3, 4, 5, 6\}$  – the number of the highest degree coefficients of  $R_n(x)$
  - Step 1.* Express the power sums  $p_{k-1}(R_n)$  and  $p_k(R_n)$  in terms of  $\{b_i\}_{i=1}^k$
  - Step 2.* Find coefficients  $\{b_i\}_{i=1}^k$  in terms of  $n$  and  $\alpha$  using Proposition 2
  - Step 3.* Find a proper value  $\sigma$  for parameter  $s$  in  $p_k - cn(n + s(\alpha + 1))p_{k-1}$ , where  $c$  is the coefficient of  $n^2$  in the quotient  $p_k/p_{k-1}$
  - Step 4.* Represent the numerator of  $f = p_k - cn(n + \sigma(\alpha + 1))p_{k-1}$  in powers of  $n$  and  $(\alpha + 1)$
  - Step 5.* Estimate from below the expression  $f$  to prove that  $f \geq 0$
- 

*Step 1:* Let  $\{x_i\}_{i=1}^n$  be all the zeros of the polynomial  $R_n(x)$  from (2.7). In order to express a power sum  $p_r = \sum_{i=1}^r x_i^r$ ,  $1 \leq r \leq n$ , by  $\{b_i\}_{i=1}^r$ , we apply the direct formula

$$p_r = \begin{vmatrix} b_1 & 1 & 0 & \dots & 0 \\ 2b_2 & b_1 & 1 & \dots & 0 \\ 3b_3 & b_2 & b_1 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots \\ rb_r & b_{r-1} & b_{r-2} & \dots & b_1 \end{vmatrix}$$

which easily follows from the Newton identities (2.11).

Below is the code of the programme and the results for  $k = 1, \dots, 6$ :

```

k = 6;
Do[p_k = Det[Table[Which[j == 1, i b_i, 1 < j ≤ i, b_{i+1-j}, j = i+1, 1, j > i+1, 0], {i, κ}, {j, κ}]];
Print[Subscript["p", κ], "=", TraditionalForm[p_k]], {κ, k}]

p_1 = b_1
p_2 = b_1^2 - 2 b_2
p_3 = b_1^3 - 3 b_2 b_1 + 3 b_3
p_4 = b_1^4 - 4 b_2 b_1^2 + 4 b_3 b_1 + 2 b_2^2 - 4 b_4
p_5 = b_1^5 - 5 b_2 b_1^3 + 5 b_3 b_1^2 + 5 b_2^2 b_1 - 5 b_4 b_1 - 5 b_2 b_3 + 5 b_5
p_6 = b_1^6 - 6 b_2 b_1^4 + 6 b_3 b_1^3 + 9 b_2^2 b_1^2 - 6 b_4 b_1^2 - 12 b_2 b_3 b_1 + 6 b_5 b_1 - 2 b_2^3 + 3 b_3^2 + 6 b_2 b_4 - 6 b_6

```

Step 2: We find coefficients  $\{b_i\}_{i=1}^k$  of the polynomial  $R_n(x)$  using Proposition 2. The source and the results for  $k = 1, \dots, 6$  follow below:

```

k = 6;
fb[κ_, n_] :=
If[κ == 1, Sum[FullSimplify[RSolveValue[{ru[q + 1] == (ru[q] + 1) (q + 1) / (q + 1 + α), ru[1] == 1 / (α + 1)}, ru[q], q]], {q, 1, n}],
Sum[Simplify[RSolveValue[{rv[q + 1] == (rv[q] + fb[κ - 1, q]) (q + 1) / (q + 1 + α), rv[1] == 0}, rv[q], q]], {q, 1, n}]]
Do[If[κ == 1, b_κ = fb[κ, n],
b_κ = Factor[Part[FactorTermsList[Numerator[fb[κ, n]], α], 2]] *
Collect[Part[FactorTermsList[Numerator[fb[κ, n]], α], 3], n, FullSimplify] / Denominator[fb[κ, n]]];
Print[Subscript["b", κ], "=", TraditionalForm[b_κ]], {κ, 1, k}]

b_1 = n (n + 1) / (2 (α + 1))
b_2 = (n - 1) n (n + 1) (3 n (α + 2) + 2 (α + 6)) / (24 (α + 1) (α + 2) (α + 3))
b_3 = (n - 2) (n - 1) n (n + 1) (5 (α + 2) (α + 4) n^2 + (α (5 α + 86) + 200) n + 12 (α + 20)) / (240 (α + 1) (α + 2) (α + 3) (α + 4) (α + 5))
b_4 = ((n - 3) (n - 2) (n - 1) n (n + 1) (105 (α + 2) (α + 4) (α + 6) n^3 + 3 (α (7 α (5 α + 204) + 9316) + 15 120) n^2 + (131 040 - 2 α (7 α (5 α + 44) - 17 244)) n - 8 (α (7 α (α + 28) + 2244) - 15 120))) / (40 320 (α + 1) (α + 2) (α + 3) (α + 4) (α + 5) (α + 6) (α + 7))
b_5 = ((n - 4) (n - 3) (n - 2) (n - 1) n (n + 1) (21 (α + 2) (α + 4) (α + 6) (α + 8) n^4 + 2 (α (α (7 α (α + 108) + 9956) + 42 928) + 56 448) n^3 + (α (α (17 988 - 7 α (7 α + 212)) + 248 496) + 572 544) n^2 + (1 241 856 - 2 α (α (21 α + 1096) + 26 468) - 34 832)) n - 240 (α (α (α + 38) + 1528) - 4032))) / (80 640 (α + 1) (α + 2) (α + 3) (α + 4) (α + 5) (α + 6) (α + 7) (α + 8) (α + 9))
b_6 = ((n - 5) (n - 4) (n - 3) (n - 2) (n - 1) n (n + 1) (3465 (α + 2) (α + 4) (α + 6) (α + 8) (α + 10) n^5 + 360 (13 α (11 α (α (7 α + 164) + 1348) + 49 936) + 739 200) n^4 + 9 (α (131 884 640 - 11 α (35 α (5 α + 278) - 10 644) - 1 805 704)) + 229 152 000) n^3 - 8 (α (11 α (α (5 α (28 α + 2685) + 620 812) + 2 759 292) - 220 067 280) - 964 656 000) n^2 + 44 (α (α (α (5 α (35 α + 1014) - 37 756) - 20 283 336) - 53 575 200) + 315 705 600) n + 96 (α (11 α (α (5 α (α + 66) + 7714) + 237 564) - 62 191 440) + 99 792 000))) / (159 667 200 (α + 1) (α + 2) (α + 3) (α + 4) (α + 5) (α + 6) (α + 7) (α + 8) (α + 9) (α + 10) (α + 11))

```

Step 3: The quotient  $p_k/p_{k-1}$  is a quadratic polynomial in  $n$ , and we denote by  $c$  its leading coefficient.

The goal of this step is to find a proper value (say  $\sigma$ ) for parameter  $s$  in the expression

$$f_s = p_k - cn(n + s(\alpha + 1))p_{k-1},$$

such that  $f_\sigma \geq 0$  for all admissible  $\alpha$  and  $n$ . For a fixed  $k$  quantity  $f_s$  depends on  $\alpha$ ,  $n$  and  $s$ . It is a polynomial of degree  $2k - 1$  in  $n$  and a rational function in  $\alpha$ . Let us write the numerator of  $f_s$  in the form

$$\sum_{i=1}^{2k-1} \sum_{j=0}^d \mu_{i,j}(s)(\alpha + 1)^{d-j} n^{2k-i}.$$

The highest order coefficients in  $\sum_j \mu_{i,j}(s)(\alpha + 1)^{d-j}$  are linear functions in  $s$  of the form  $A_i - B_i s$ , with  $A_i > 0$  and  $B_i > 0$ . We denote their zeros by  $s_i$  for each  $i$  and set  $\sigma = \min_i s_i$ . Since we seek estimates valid for all  $\alpha > -1$ , our choice of  $\sigma$  guarantee that for  $\alpha$  sufficiently large the inequality  $\sum_j \mu_{i,j}(s)(\alpha + 1)^{d-j} > 0$  holds true.

The code is as follows:

```
r = PolynomialQuotient[pk, p[k-1], n];
c = Factor[Coefficient[r, n, 2]];
fs = pk - c n (n + s (alpha + 1)) p[k-1];
numfs = Numerator[Together[Apart[fs, alpha]]]
Do[gs = Factor[Coefficient[numfs, n, i]];
  num = Normal[Series[gs, {alpha, -1, Exponent[gs, alpha]}]];
  sols = Solve[Coefficient[num, alpha, Exponent[gs, alpha]] = 0, s, Reals];
  ss[i] = s /. Flatten[sols], {i, 2 k - 1, 1, -1}];
sigma = Min[Table[ss[i], {i, 2, 2 k - 1}]];
```

Table 1 gives results for the optimal values of  $c$  and  $\sigma$  for  $k = 3, 4, 5, 6$ .

Table 1: The optimal values of  $c$  and  $\sigma$  in the lower bounds for  $c_n^2(\alpha)$ .

$k$	$c$	$\sigma$
3	$\frac{2}{(\alpha + 1)(\alpha + 5)}$	$\frac{3}{8}$
4	$\frac{5\alpha + 17}{2(\alpha + 1)(\alpha + 3)(\alpha + 7)}$	$\frac{8}{25}$
5	$\frac{2(7\alpha + 31)}{(\alpha + 1)(\alpha + 9)(5\alpha + 17)}$	$\frac{25}{84}$
6	$\frac{21\alpha^3 + 299\alpha^2 + 1391\alpha + 2073}{(\alpha + 1)(\alpha + 3)(\alpha + 5)(\alpha + 11)(7\alpha + 31)}$	$\frac{2}{7}$

Step 4: We set

$$f = p_k - cn(n + \sigma(\alpha + 1))p_{k-1} =: \frac{\varphi(n, \alpha)}{\psi(\alpha)}$$

with  $c$  and  $\sigma$  determined in Step 3. Here,  $\varphi(n, \alpha)$  is a bivariate polynomial in  $n$  and  $\alpha$ , and  $\psi(\alpha)$  is a polynomial in  $\alpha$ . More precisely,  $\varphi(n, \alpha)$  has degree  $2k-1$  in  $n$ , and degree  $d$  in  $\alpha$  which our programme calculates for each fixed  $k$ .

Note that  $\psi(\alpha) > 0$  for  $\alpha > -1$  since it is a product of powers of  $\alpha + j$ ,  $j \geq 1$  and multipliers  $A\alpha + B$ ,  $0 < A < B$ . Therefore,  $\text{sign } f = \text{sign } \varphi$ .

We expand  $\varphi(n, \alpha)$  in the form

$$\varphi(n, \alpha) = \sum_{i=1}^{2k-1} \sum_{j=0}^d \mu_{i,j} (\alpha + 1)^{d-j} n^{2k-i} = \begin{pmatrix} n^{2k-1} \\ n^{2k-2} \\ \vdots \\ n \end{pmatrix}^{\top} \mathbf{M} \begin{pmatrix} (\alpha + 1)^d \\ (\alpha + 1)^{d-1} \\ \vdots \\ 1 \end{pmatrix},$$

where  $\mathbf{M} = (\mu_{i,j})_{i=1,j=0}^{2k-1,d}$  and all entries  $\mu_{i,j}$  are integer numbers.

The source for computation of the matrix  $\mathbf{M}$  is listed below.

```
f = p_k - c n (n + sigma (alpha + 1)) p_{k-i};
psi = Numerator[Together[Apart[f, alpha]]];
psi = Denominator[Together[Apart[f, alpha]]];
Do[g = Factor[Coefficient[psi, n, i]]; dag[i] = Exponent[g, alpha], {i, 2 k - 1, 1, -1}]
d = Max[Table[dag[i], {i, 1, 2 k - 1}]] + 1;
mu = ConstantArray[0, {2 k - 1, d}];
Do[g = Factor[Coefficient[psi, n, i]];
  Table[mu[[2 k - i, d - j]] = SeriesCoefficient[Series[g, {alpha, -1, dag[i]}], j], {j, 0, dag[i]}],
  {i, 2 k - 1, 1, -1}];
```

If  $\mu_{i,j} \geq 0$  for all  $i, j$ , then  $\varphi(n, \alpha) \geq 0$  and  $f \geq 0$  for all  $\alpha > -1$  and  $n \geq k$ . In a case some of coefficients  $\mu_{i,j} < 0$  we apply the next step of the algorithm.

The results for  $k = 3, 4, 5, 6$  are given together with the estimates from Step 5.

*Step 5:* If there are coefficients  $\mu_{i,j} < 0$  we need additional arguments to verify that  $f \geq 0$  for all  $\alpha > -1$  and  $n \geq k$ . We bring into use a new  $(2k-1) \times (d+1)$  matrix  $\mathbf{\Lambda}$  which elements we put initially  $\lambda_{i,j} := \mu_{i,j}$ , for  $i = 1, \dots, 2k-1$  and  $j = 0, \dots, d$ .

The procedure described below checks recursively all coefficients  $\lambda_{i,j}$  and makes the corresponding estimations. We need not introduce a new matrix after each iteration, but only replace a pair of elements in a column of  $\mathbf{\Lambda}$  with new entries in such a manner that the value of the function

$$\Phi(\mathbf{\Lambda}) = \sum_{i=1}^{2k-1} \sum_{j=0}^d \lambda_{i,j} (\alpha + 1)^{d-j} n^{2k-i} = \begin{pmatrix} n^{2k-1} \\ n^{2k-2} \\ \vdots \\ n \end{pmatrix}^{\top} \mathbf{\Lambda} \begin{pmatrix} (\alpha + 1)^d \\ (\alpha + 1)^{d-1} \\ \vdots \\ 1 \end{pmatrix}$$

decreases. At the end of the procedure we get a matrix  $\mathbf{\Lambda}$  satisfying  $\mathbf{0} \leq \mathbf{\Lambda} \leq \mathbf{M}$  (in the sense that  $0 \leq \lambda_{i,j} \leq \mu_{i,j}$  for all  $i, j$ ) and therefore

$$0 \leq \Phi(\mathbf{\Lambda}) \leq \Phi(\mathbf{M}) = \varphi(n, \alpha).$$

Suppose that  $\lambda_{i,j} < 0$  for some pair of indices  $i, j$ . Then we set

$$h := \min\{i - \eta : \lambda_{\eta,j} > 0, 1 \leq \eta \leq i - 1\} \quad \text{and} \quad \delta := \frac{\lambda_{i,j}}{k^{i-h}} \quad (\delta < 0).$$

If  $\lambda_{h,j} + \delta \geq 0$ , for  $n \geq k$  we have

$$\begin{aligned} (\lambda_{h,j} + \delta)n^{2k-h} + 0n^{2k-i} &= \left(\lambda_{h,j} + \frac{\lambda_{i,j}}{k^{i-h}}\right)n^{2k-h} = \lambda_{h,j}n^{2k-h} + \lambda_{i,j}\frac{n^{2k-h}}{k^{i-h}} \\ &\leq \lambda_{h,j}n^{2k-h} + \lambda_{i,j}\frac{n^{2k-h}}{n^{i-h}} = \lambda_{h,j}n^{2k-h} + \lambda_{i,j}n^{2k-i}. \end{aligned}$$

Otherwise, if  $\lambda_{h,j} + \delta < 0$ , for  $n \geq k$  we have

$$\begin{aligned} 0n^{2k-h} + (\lambda_{h,j}k^{i-h} + \lambda_{i,j})n^{2k-i} &= \lambda_{h,j}n^{2k-i}k^{i-h} + \lambda_{i,j}n^{2k-i} \\ &\leq \lambda_{h,j}n^{2k-i}n^{i-h} + \lambda_{i,j}n^{2k-i} \\ &\leq \lambda_{h,j}n^{2k-h} + \lambda_{i,j}n^{2k-i}. \end{aligned}$$

So, replacing only two elements in  $\mathbf{\Lambda}$ ,

$$\begin{cases} \lambda_{h,j} := \lambda_{h,j} + \lfloor \delta \rfloor & \text{and } \lambda_{i,j} := 0, & \text{if } \lambda_{h,j} + \delta \geq 0, \\ \lambda_{i,j} := \lambda_{h,j}k^{i-h} + \lambda_{i,j} & \text{and } \lambda_{h,j} := 0, & \text{otherwise,} \end{cases}$$

we obtain that

$$\lambda_{h,j}(\alpha + 1)^{d+1-j}n^{2k-h} + \lambda_{i,j}(\alpha + 1)^{d+1-j}n^{2k-i}$$

decreases for the new values of  $\lambda_{h,j}$  and  $\lambda_{i,j}$ , and hence  $\Phi(\mathbf{\Lambda})$  also decreases.

Applying recursively the above iteration process for  $i = 2k - 1, 2k - 2, \dots, 1$  and  $j = 0, 1, \dots, d$  we finally obtain a matrix  $\mathbf{\Lambda}$  satisfying  $\mathbf{0} \leq \mathbf{\Lambda} \leq \mathbf{M}$ . Then  $\varphi(n, \alpha) \geq 0$ ,  $f \geq 0$  and therefore

$$c_n^2(\alpha) \geq \frac{p_k}{p_{k-1}} \geq cn(n + \sigma(\alpha + 1))$$

for the optimal  $c$  and  $\sigma$  evaluated in Step 3. For  $k = 3, 4, 5, 6$  we obtain estimates (1.10), (1.12), (1.14), and (1.16), respectively.

The following source implements the procedure described in Step 5.

```

λ = μ;
For[i = 2 k - 1, i > 1, i--]
  For[j = 1, j ≤ d, j++, If[λ[[i, j]] ≥ 0, Continue[]];
  h = i - First[FirstPosition[Positive[λ[[i - 1 ;; 1 ;; -1, j]], True]];
  δ = λ[[i, j]]/(k^(i - h));
  If[λ[[h, j]] + δ ≥ 0, λ[[h, j]] = λ[[h, j]] + Floor[δ]; λ[[i, j]] = 0,
  λ[[i, j]] = λ[[h, j]] * k^(i - h) + λ[[i, j]]; λ[[h, j]] = 0; i = i + 1]]
Print["Λ = ", MatrixForm[λ]]
Print["M = ", MatrixForm[μ]]

```

Next, we give matrices  $\mathbf{M}$  from Step 4 and  $\mathbf{\Lambda}$  from Step 5 obtained with *Mathematica*.

Case  $k = 3$ :

This partial case needs a special attention as we have to assume strict inequality  $n > k$ , i.e.,  $n \geq 4$ , to obtain estimate (1.10). This causes a minor modification in Step 5 of Algorithm 1, namely, replacement of  $k^{i-h}$  with  $(k+1)^{i-h}$ . Namely, we determine  $\delta := \lambda_{i,j}/(k+1)^{i-h}$  and set

$$\begin{cases} \lambda_{h,j} := \lambda_{h,j} + \lfloor \delta \rfloor & \text{and } \lambda_{i,j} := 0, \quad \text{if } \lambda_{h,j} + \delta \geq 0, \\ \lambda_{i,j} := \lambda_{h,j} (k+1)^{i-h} + \lambda_{i,j} & \text{and } \lambda_{h,j} := 0, \quad \text{otherwise.} \end{cases}$$

Matrices  $\mathbf{M}$  and  $\mathbf{\Lambda}$  in this case are

$$\mathbf{\Lambda} = \begin{pmatrix} 0 & 4 & -4 & 225 & 360 \\ 0 & 0 & 390 & 510 & 720 \\ 15 & 155 & 205 & 1185 & 360 \\ 15 & 270 & 495 & 900 & 0 \\ 0 & 36 & 684 & 0 & 0 \end{pmatrix} \quad \mathbf{M} = \begin{pmatrix} 0 & 19 & -4 & 225 & 360 \\ 0 & -60 & 390 & 510 & 720 \\ 15 & 155 & 205 & 1185 & 360 \\ 15 & 270 & 495 & 900 & 0 \\ 0 & 36 & 684 & 0 & 0 \end{pmatrix}.$$

Although there is a negative element of  $\mathbf{\Lambda}$ , from  $4(\alpha+1)^2 - 4(\alpha+1) + 225 \geq 0$  for all  $\alpha > -1$  we conclude that  $4(\alpha+1)^3 - 4(\alpha+1)^2 + 225(\alpha+1) + 360 > 0$  and consequently  $\Phi(\mathbf{\Lambda}) \geq 0$  for  $n \geq 4$ .

By a direct verification one can see that inequality (1.10) holds also in the case  $n = k = 3$ .

Case  $k = 4$ :

$$\mathbf{\Lambda} = \begin{pmatrix} 0 & 0 & 10200 & 72480 & 323700 & 1413060 & 3602340 & 4340700 & 1890000 \\ 0 & 4882 & 30891 & 359695 & 2625259 & 7966210 & 13275570 & 12707100 & 5670000 \\ 0 & 0 & 229110 & 1642830 & 6282570 & 16699200 & 24837120 & 18692100 & 5670000 \\ 2100 & 46515 & 120645 & 2404465 & 10159765 & 20026720 & 25810890 & 16625700 & 1890000 \\ 2756 & 106120 & 876330 & 2582090 & 7616630 & 17567550 & 18060000 & 6300000 & 0 \\ 0 & 11060 & 662604 & 2653840 & 6215776 & 11121880 & 7413000 & 0 & 0 \\ 0 & 0 & 0 & 1120600 & 4777900 & 3435000 & 0 & 0 & 0 \end{pmatrix}$$

$$\mathbf{M} = \begin{pmatrix} 0 & 0 & 10200 & 72480 & 323700 & 1413060 & 3602340 & 4340700 & 1890000 \\ 0 & 8715 & 30891 & 359695 & 2625259 & 7966210 & 13275570 & 12707100 & 5670000 \\ 0 & -15330 & 229110 & 1642830 & 6282570 & 16699200 & 24837120 & 18692100 & 5670000 \\ 2100 & 46515 & 120645 & 2404465 & 10159765 & 20026720 & 25810890 & 16625700 & 1890000 \\ 2800 & 106120 & 876330 & 2582090 & 7616630 & 17567550 & 18060000 & 6300000 & 0 \\ 0 & 15960 & 722904 & 2653840 & 6215776 & 11121880 & 7413000 & 0 & 0 \\ -700 & -19600 & -241200 & 1120600 & 4777900 & 3435000 & 0 & 0 & 0 \end{pmatrix}$$

Case  $k = 5$ :

$$\mathbf{\Lambda} = \begin{pmatrix} 0 & 0 & 0 & 64925 & 1064665 & 8138830 & 43256150 & 172898565 & 474925185 & 805850640 & 734423760 & 266716800 \\ 0 & 0 & 91665 & 1204470 & 9699090 & 71280390 & 373661895 & 1241223900 & 2610599670 & 3473555400 & 2804336640 & 1066867200 \\ 0 & 19824 & 130578 & 3408188 & 48487642 & 313463920 & 1271550350 & 3522779568 & 6544523790 & 7686433440 & 5117787360 & 1600300800 \\ 0 & 0 & 1451982 & 16288020 & 114900450 & 672910770 & 2546690160 & 6152610870 & 9859721760 & 10218685680 & 5871579840 & 1066867200 \\ 3675 & 128835 & 0 & 24490445 & 226233910 & 991504675 & 3153540110 & 7169071245 & 10438959825 & 9013742640 & 3935025360 & 266716800 \\ 6027 & 381850 & 6416795 & 22404550 & 169885205 & 1005110890 & 2985302145 & 5744010510 & 7716554370 & 5584488840 & 1111320000 & 0 \\ 0 & 52297 & 5062484 & 58263912 & 213196158 & 589342950 & 1804792500 & 3787471002 & 4038237000 & 1770703200 & 0 & 0 \\ 0 & 0 & 0 & 15084950 & 144208510 & 409403975 & 1057769610 & 1931913900 & 1309770000 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 256255650 & 690284700 & 417538800 & 0 & 0 & 0 & 0 \end{pmatrix}$$

$$\mathbf{M} = \begin{pmatrix} 0 & 0 & 0 & 64925 & 1064665 & 8138830 & 43256150 & 172898565 & 474925185 & 805850640 & 734423760 & 266716800 \\ 0 & 0 & 91665 & 1204470 & 9699090 & 71280390 & 373661895 & 1241223900 & 2610599670 & 3473555400 & 2804336640 & 1066867200 \\ 0 & 27804 & 130578 & 3408188 & 48487642 & 313463920 & 1271550350 & 3522779568 & 6544523790 & 7686433440 & 5117787360 & 1600300800 \\ 0 & -39900 & 1500030 & 16288020 & 114900450 & 672910770 & 2546690160 & 6152610870 & 9859721760 & 10218685680 & 5871579840 & 1066867200 \\ 3675 & 128835 & -240240 & 24490445 & 226233910 & 991504675 & 3153540110 & 7169071245 & 10438959825 & 9013742640 & 3935025360 & 266716800 \\ 6125 & 381850 & 6416795 & 22404550 & 169885205 & 1005110890 & 2985302145 & 5744010510 & 7716554370 & 5584488840 & 1111320000 & 0 \\ 0 & 77616 & 5699022 & 58263912 & 213196158 & 589342950 & 1804792500 & 3787471002 & 4038237000 & 1770703200 & 0 & 0 \\ -2450 & -123445 & -3055430 & 20292530 & 152590030 & 409403975 & 1057769610 & 1931913900 & 1309770000 & 0 & 0 & 0 \\ 0 & -15750 & -636300 & -26037900 & -41907600 & 256255650 & 690284700 & 417538800 & 0 & 0 & 0 & 0 \end{pmatrix}$$

Case  $k = 6$ :

$$\begin{aligned}
 \Lambda^T &= \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 48510 & 95223 & 0 & 0 & 0 & 16170 \\ 0 & 0 & 0 & 425810 & 0 & 2817045 & 9741270 & 1348462 & 0 & 0 & 1252020 \\ 0 & 0 & 3476550 & 6110115 & 48434732 & 0 & 336258384 & 218861747 & 0 & 0 & 38848656 \\ 0 & 6055665 & 95465370 & 190273710 & 1221447150 & 1171139970 & 2726237052 & 7298343195 & 0 & 0 & 1158647028 \\ 3128160 & 204531195 & 1480047030 & 5336244870 & 15771654360 & 32618391960 & 21628131756 & 73442566505 & 29020437224 & 0 & 0 \\ 116263280 & 3318028175 & 18873326010 & 77596724865 & 174458095350 & 356484794820 & 298526146072 & 392659895320 & 419332019003 & 0 & 0 \\ 1988081620 & 35746404925 & 197029544250 & 726747979015 & 1531387171180 & 2562636437130 & 275434379016 & 1907270574440 & 2403530867430 & 384166681454 & 0 \\ 2110299620 & 29036961329 & 155826940290 & 496082042100 & 10157416978170 & 1439765320512 & 15830016304564 & 10335346465675 & 8454107203080 & 372137898370 & 0 \\ 15752193380 & 1842856573327 & 9151953918030 & 2582662738780 & 49078270584420 & 64463871381756 & 64790433176084 & 46215397662665 & 25413887653770 & 13866272170542 & 3624993260826 \\ 879576036500 & 909732993521 & 40326294432270 & 103410904320900 & 179681190528840 & 232720508502183 & 206003553429058 & 148200002432020 & 74515561079190 & 38329760467746 & 21352330210512 \\ 37689721407020 & 34425402760287 & 13493782918400 & 317406347163180 & 506548245985320 & 592781349468231 & 516092578758680 & 349590783269990 & 182484813042840 & 861307227278092 & 36302262824520 \\ 12408373123020 & 9838568531450 & 344232626910300 & 742720322283380 & 109886468687920 & 1195101212250330 & 98299659161584 & 622626376181040 & 315159824447160 & 127817022168000 & 2759433993216 \\ 30888195143400 & 211221314490186 & 667161153364860 & 1314049020225480 & 180585150980850 & 181528345062778 & 1389893720693940 & 808485195464100 & 350770307406760 & 103510976206176 & 7959911420160 \\ 56418683248620 & 33395069661060 & 961433219937960 & 1730658737031840 & 2184547015159740 & 2011470759046980 & 1371614133582000 & 691021013177880 & 227021138467200 & 3392761477760 & 0 \\ 72303760123560 & 380597496158880 & 996948089953280 & 164078227386560 & 184198023153040 & 1520944502505120 & 876543356468320 & 330199345808640 & 64746646272000 & 0 & 0 \\ 60744201708960 & 29769661581600 & 70488552078400 & 1045548136987200 & 1036553459911200 & 695810390758560 & 300533746867200 & 6368871238400 & 0 & 0 & 0 \\ 29689237670400 & 143165195712000 & 307454361984000 & 39195024424000 & 317437952832000 & 151152068390400 & 31685955840000 & 0 & 0 & 0 & 0 \\ 6337191168000 & 31685955840000 & 63371911680000 & 63371911680000 & 31685955840000 & 63371911680000 & 31685955840000 & 63371911680000 & 63371911680000 & 63371911680000 & 63371911680000 \end{pmatrix} \\
 \Lambda^T &= \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 48510 & 97020 & 0 & -64680 & 0 & 16170 \\ 0 & 0 & 0 & 54005 & -709170 & 2817045 & 9741270 & 2279970 & -5453910 & -810810 & 1252020 \\ 0 & 0 & 3476550 & 6110115 & 5141590 & -1788745 & 336258384 & 22910660 & -23328490 & -5087310 & 38848656 \\ 0 & 6055665 & 95465370 & 190273710 & 1221447150 & 1171139970 & 2726237052 & 752825695 & -62753080 & -431601520 & 1158647028 \\ 3128160 & 204531195 & 1480047030 & 5336244870 & 15771654360 & 32618391960 & 21628131756 & 73442566505 & 36182631870 & 4253399320 & -5517429876 \\ 116263280 & 3318028175 & 18873326010 & 77596724865 & 174458095350 & 356484794820 & 298526146072 & 392659895320 & 439184120760 & 47653879280 & -180752387572 \\ 1988081620 & 35746404925 & 197029544250 & 726747979015 & 1531387171180 & 2562636437130 & 275434379016 & 1907270574440 & 2403530867430 & 656697131530 & -1107182700456 \\ 2110299620 & 29036961329 & 155826940290 & 496082042100 & 10157416978170 & 1439765320512 & 15830016304564 & 10335346465675 & 8454107203080 & 4052140160904 & -1984507575204 \\ 15752193380 & 1842856573327 & 9151953918030 & 2582662738780 & 49078270584420 & 64463871381756 & 64790433176084 & 46215397662665 & 25413887653770 & 13866272170542 & 3624993260826 \\ 879576036500 & 909732993521 & 40326294432270 & 103410904320900 & 179681190528840 & 232720508502183 & 206003553429058 & 148200002432020 & 74515561079190 & 38329760467746 & 21352330210512 \\ 37689721407020 & 34425402760287 & 13493782918400 & 317406347163180 & 506548245985320 & 592781349468231 & 516092578758680 & 349590783269990 & 182484813042840 & 861307227278092 & 36302262824520 \\ 12408373123020 & 9838568531450 & 344232626910300 & 742720322283380 & 109886468687920 & 1195101212250330 & 98299659161584 & 622626376181040 & 315159824447160 & 127817022168000 & 2759433993216 \\ 30888195143400 & 211221314490186 & 667161153364860 & 1314049020225480 & 180585150980850 & 181528345062778 & 1389893720693940 & 808485195464100 & 350770307406760 & 103510976206176 & 7959911420160 \\ 56418683248620 & 33395069661060 & 961433219937960 & 1730658737031840 & 2184547015159740 & 2011470759046980 & 1371614133582000 & 691021013177880 & 227021138467200 & 3392761477760 & 0 \\ 72303760123560 & 380597496158880 & 996948089953280 & 164078227386560 & 184198023153040 & 1520944502505120 & 876543356468320 & 330199345808640 & 64746646272000 & 0 & 0 \\ 60744201708960 & 29769661581600 & 70488552078400 & 1045548136987200 & 1036553459911200 & 695810390758560 & 300533746867200 & 6368871238400 & 0 & 0 & 0 \\ 29689237670400 & 143165195712000 & 307454361984000 & 39195024424000 & 317437952832000 & 151152068390400 & 31685955840000 & 0 & 0 & 0 & 0 \\ 6337191168000 & 31685955840000 & 63371911680000 & 63371911680000 & 31685955840000 & 63371911680000 & 31685955840000 & 63371911680000 & 63371911680000 & 63371911680000 & 63371911680000 \end{pmatrix}
 \end{aligned}$$

### 3.2. UPPER BOUNDS FOR $c_n(\alpha)$

We apply Proposition 3 to estimate the largest zero  $x_n = c_n^2(\alpha)$  of the polynomial  $R_n(x)$  from above,

$$x_n \leq u_k(R_n) = p_k(R_n)^{1/k}, \quad k = 3, 4, 5, 6.$$

Then with the assistance of computer algebra we obtain a further estimation of the form

$$u_k(R_n) \leq c^{1/k} (n+1)(n+\sigma(\alpha+1)),$$

with the optimal (i.e., the smallest possible) constants  $c = c(k)$  and  $\sigma = \sigma(k)$ .

The algorithm is analogous to Algorithm 1, and the code has only a few differences which are specified later.

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**Algorithm 2** Estimating  $c_n(\alpha)$  from above

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- Input:*  $k \in \{3, 4, 5, 6\}$  – the number of the highest degree coefficients of  $R_n(x)$
  - Step 1.* Express the power sum  $p_k(R_n)$  in terms of  $\{b_i\}_{i=1}^k$
  - Step 2.* Find  $\{b_i\}_{i=1}^k$  in terms of  $n$  and  $\alpha$  using Proposition 2
  - Step 3.* Find a proper value  $\sigma$  for parameter  $s$  in the expression  $c(n+1)^k(n+s(\alpha+1))^k - p_k$ , where  $c$  is the coefficient of  $n^{2k}$  in  $p_k$
  - Step 4.* Represent the numerator of  $f = c(n+1)^k(n+s(\alpha+1))^k - p_k$  in powers of  $n$  and  $(\alpha+1)$
  - Step 5.* Estimate from below the expression  $f$  to prove that  $f \geq 0$
- 

Step 1: The same as in Algorithm 1.

Step 2: Identical to that in Algorithm 1.



Step 3: The only differences with Algorithm 1 are that we set  $c$  to be the coefficient of  $n^{2k}$  in  $p_k$  and

$$f_s = c(n+1)^k(n+s(\alpha+1))^k - p_k.$$

The highest order coefficients in  $\sum_j \mu_{i,j}(s)(\alpha+1)^{d-j}$  are functions in  $s$  of the form  $A_i s^\nu - B_i$ , with  $A_i > 0$  and  $B_i \geq 0$ . We denote their non-negative zeros by  $s_i$  for each  $i$  and choose  $\sigma = \max_i s_i$ .

The results for  $k = 3, 4, 5, 6$  obtained by symbolic computations are given in Table 2.

Table 2: The optimal values of  $c$  and  $\sigma$  in the upper bounds for  $c_n^2(\alpha)$ .

$k$	$c$	$\sigma$
3	$\frac{1}{(\alpha+1)^3(\alpha+3)(\alpha+5)}$	$\frac{2}{5}$
4	$\frac{5\alpha+17}{2(\alpha+1)^4(\alpha+3)^2(\alpha+5)(\alpha+7)}$	$\frac{3}{7}$
5	$\frac{(7\alpha+31)}{(\alpha+1)^5(\alpha+3)^2(\alpha+5)(\alpha+7)(\alpha+9)}$	$\frac{4}{9}$
6	$\frac{21\alpha^3+299\alpha^2+1391\alpha+2073}{(\alpha+1)^6(\alpha+3)^3(\alpha+5)^2(\alpha+7)(\alpha+9)(\alpha+11)}$	$\frac{5}{11}$

Step 4: With  $c$  and  $\sigma$  determined in the previous Step 3 we set

$$f = c(n+1)^k(n+\sigma(\alpha+1))^k - p_k =: \frac{\varphi(n, \alpha)}{\psi(\alpha)}.$$

The rest of the source has no difference with Step 4 of Algorithm 1.

Step 5: The same as in Algorithm 1. Using the same recursive procedure we find a matrix  $\mathbf{\Lambda}$  satisfying  $\mathbf{0} \leq \mathbf{\Lambda} \leq \mathbf{M}$ . Then  $\varphi(n, \alpha) \geq 0$ ,  $f \geq 0$  and therefore

$$c_n^{2k}(\alpha) \leq p_k \leq c(n+1)^k(n+\sigma(\alpha+1))^k$$

for the corresponding  $c$  and  $\sigma$  evaluated in Step 3. For  $k = 3, 4, 5, 6$  we obtain estimations (1.11), (1.13), (1.15), and (1.17), respectively.

The matrices  $\mathbf{M}$  from Step 4 and  $\mathbf{\Lambda}$  from Step 5 obtained with *Mathematica* are given below.

Case  $k = 3$ :

$$\mathbf{\Lambda} = \begin{pmatrix} 0 & 0 & 0 & 1500 & 3300 \\ 0 & 115 & 1885 & 4170 & 4233 \\ 32 & 598 & 3026 & 6360 & 0 \\ 96 & 979 & 2143 & 850 & 0 \\ 96 & 624 & 1098 & 0 & 0 \end{pmatrix} \quad \mathbf{M} = \begin{pmatrix} 0 & 0 & 0 & 1500 & 3300 \\ 0 & 115 & 1885 & 4170 & 4650 \\ 32 & 598 & 3026 & 6360 & -600 \\ 96 & 979 & 2143 & 1560 & -1950 \\ 96 & 624 & 1098 & -2130 & 0 \end{pmatrix}$$

Case  $k = 4$ :

$$\Lambda = \begin{pmatrix} 0 & 0 & 0 & 0 & 905520 & 8808240 & 29717520 & 41571600 & 19756800 \\ 0 & 0 & 54390 & 2038890 & 16676660 & 60285680 & 115770830 & 117031110 & 48774600 \\ 0 & 42294 & 1237572 & 10966494 & 52723608 & 141477042 & 198565500 & 127823850 & 24194362 \\ 6075 & 266115 & 3694950 & 25364010 & 85166735 & 157047575 & 154257320 & 46893642 & 0 \\ 24300 & 617510 & 5700800 & 26734470 & 72437020 & 97039330 & 34815501 & 0 & 0 \\ 36450 & 678780 & 4979940 & 16392810 & 28823750 & 17907835 & 0 & 0 & 0 \\ 24300 & 360421 & 2131108 & 6792156 & 5246162 & 0 & 0 & 0 & 0 \end{pmatrix}$$

$$M = \begin{pmatrix} 0 & 0 & 0 & 0 & 905520 & 8808240 & 29717520 & 41571600 & 19756800 \\ 0 & 0 & 54390 & 2038890 & 16676660 & 60285680 & 115770830 & 117031110 & 48774600 \\ 0 & 42294 & 1237572 & 10966494 & 52723608 & 141477042 & 198565500 & 127823850 & 27783000 \\ 6075 & 266115 & 3694950 & 25364010 & 85166735 & 157047575 & 154257320 & 52558380 & -11730600 \\ 24300 & 617510 & 5700800 & 26734470 & 72437020 & 97039330 & 38636640 & -18088350 & -10495800 \\ 36450 & 678780 & 4979940 & 16392810 & 28823750 & 20280800 & -12849340 & -18282390 & 0 \\ 24300 & 360421 & 2131108 & 6792156 & 5246162 & -9491857 & -9740850 & 0 & 0 \end{pmatrix}$$

Case  $k = 5$ :

$$\Lambda = \begin{pmatrix} 0 & 0 & 0 & 0 & 85424220 & 1436596560 & 8988832440 & 26097558480 & 34662943980 & 16203045600 \\ 0 & 0 & 0 & 4261005 & 260814330 & 3617057430 & 22550151630 & 73071107235 & 134891273160 & 134642808090 & 56710659600 \\ 0 & 0 & 5436720 & 241567920 & 3235204800 & 2246774740 & 91003127400 & 223063050420 & 312360753600 & 232393230640 & 64832182400 \\ 0 & 1982358 & 88937982 & 1392482448 & 12340605438 & 63755213760 & 194677526736 & 357163148790 & 375802372260 & 186521488020 & 12638375568 \\ 200704 & 14563010 & 340432890 & 4020858058 & 25446365294 & 99455228208 & 241336266948 & 338611016520 & 235926284580 & 44541786567 & 0 \\ 1003520 & 42390775 & 693405300 & 6004806185 & 31876009900 & 96870254355 & 175080003840 & 176585507595 & 54286938720 & 0 & 0 \\ 2007040 & 63580160 & 829630410 & 5638883530 & 22495811450 & 57112266330 & 77686343280 & 30853075478 & 0 & 0 & 0 \\ 2007040 & 52428341 & 568553244 & 3375204826 & 9950248616 & 17535199185 & 13032227178 & 0 & 0 & 0 & 0 \\ 1003520 & 22758400 & 207566490 & 998218460 & 3486984100 & 3092469120 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

$$M = \begin{pmatrix} 0 & 0 & 0 & 0 & 85424220 & 1436596560 & 8988832440 & 26097558480 & 34662943980 & 16203045600 \\ 0 & 0 & 0 & 4261005 & 260814330 & 3617057430 & 22550151630 & 73071107235 & 134891273160 & 134642808090 & 56710659600 \\ 0 & 0 & 5436720 & 241567920 & 3235204800 & 2246774740 & 91003127400 & 223063050420 & 312360753600 & 232393230640 & 64832182400 \\ 0 & 1982358 & 88937982 & 1392482448 & 12340605438 & 63755213760 & 194677526736 & 357163148790 & 375802372260 & 186521488020 & 12630345600 \\ 200704 & 14563010 & 340432890 & 4020858058 & 25446365294 & 99455228208 & 241336266948 & 338611016520 & 235926284580 & 51689001420 & -16203045600 \\ 1003520 & 42390775 & 693405300 & 6004806185 & 31876009900 & 96870254355 & 175080003840 & 176585507595 & 59214803760 & -31849915230 & -8101522800 \\ 2007040 & 63580160 & 829630410 & 5638883530 & 22495811450 & 57112266330 & 77686343280 & 32878980540 & -21278795580 & 0 & 0 \\ 2007040 & 52428341 & 568553244 & 3375204826 & 9950248616 & 17535199185 & 14090589072 & -8987585040 & -16802648100 & 0 & 0 \\ 1003520 & 22758400 & 207566490 & 998218460 & 3486984100 & 3092469120 & -5291809470 & -5709701340 & 0 & 0 & 0 \end{pmatrix}$$

Case  $k = 6$ :

$$\Lambda^T = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 17812500 & 82687800 & 286787500 & 276220000 & 286787500 & 82687800 \\ 0 & 0 & 0 & 0 & 0 & 1712831340 & 16225847190 & 6101397190 & 12162937500 & 141073384270 & 9616125000 & 3500193900 \\ 0 & 0 & 0 & 7214978925 & 129414811020 & 682431963570 & 1909641192600 & 3176038395495 & 323829069980 & 19074664930 & 68515048600 \\ 0 & 0 & 11887522200 & 49669292025 & 397563217280 & 158452811640 & 349398172095 & 4862826140345 & 4402671735710 & 242950992370 & 768368197232 \\ 0 & 675992780 & 889711262440 & 1292197935670 & 718806591680 & 21861776740580 & 407202492428 & 488239798537030 & 39180748052700 & 1986889884620 & 5393481373052 \\ 0 & 6782207080 & 238370724380 & 2041791514140 & 8480781863740 & 208333346883670 & 332266439747106 & 34888681378940 & 24097462767010 & 1108868783760 & 28043348858044 \\ 57878637760 & 28564042298220 & 3098176511479280 & 14356440238205635 & 3991199974786200 & 6987865311080820 & 81762211401327198 & 6521352883777995 & 338328730973790 & 115312720216130 & 28509977559272 \\ 498282579760 & 252414547688728 & 194679747235820 & 7809258025046420 & 165724808432320 & 233328219344424 & 25297282028181852 & 1748441876085540 & 7870665648007900 & 21129116080908 & 377409107237628 \\ 82687442447120 & 146638748071388 & 896629728204120 & 2704029272288830 & 52500679384868280 & 67526748901207324 & 8485864724419364 & 33707337050110470 & 1263780830762420 & 2517631025243518 & 100310771884846 \\ 148140314822640 & 5881144622708016 & 27906262771514000 & 7240085337919890 & 198318687198670400 & 131071181160122648 & 9600891421160308 & 4523389406272600 & 1254982652811640 & 108242769642964 & 0 \\ 1754184701373740 & 1607733127380292 & 6335138217720660 & 1404602584860210 & 1972501870281840 & 1819808232340796 & 10380671394478244 & 386368920976090 & 501172233370481 & 0 & 0 \\ 439673140007600 & 31771283186568156 & 1035792171132940 & 193938078844401000 & 22662178817967640 & 1716811248944619828 & 768783125513892720 & 14726562829687989 & 0 & 0 & 0 \\ 70706027359440 & 42142727981133820 & 11662866391620320 & 187004174964081860 & 17633830408979200 & 10063921329237760 & 26672860281531376 & 0 & 0 & 0 & 0 \\ 712526958986440 & 3786625201053360 & 8577182988763280 & 1127420403129960 & 820524248199040 & 2776642514174840 & 0 & 0 & 0 & 0 & 0 \\ 4015182786100800 & 190178878072832800 & 371893024262944800 & 3747643864777600 & 1763962690692400 & 0 & 0 & 0 & 0 & 0 & 0 \\ 95728301094400 & 452632838792000 & 7271669102544000 & 496172949856400 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

$$M^T = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 17812500 & 82687800 & 286787500 & 276220000 & 286787500 & 82687800 \\ 0 & 0 & 0 & 0 & 0 & 1712831340 & 16225847190 & 6101397190 & 12162937500 & 141073384270 & 9616125000 & 3500193900 \\ 0 & 0 & 0 & 7214978925 & 129414811020 & 682431963570 & 1909641192600 & 3176038395495 & 323829069980 & 19074664930 & 68515048600 \\ 0 & 0 & 11887522200 & 49669292025 & 397563217280 & 158452811640 & 349398172095 & 4862826140345 & 4402671735710 & 242950992370 & 768368197232 \\ 0 & 675992780 & 889711262440 & 1292197935670 & 718806591680 & 21861776740580 & 407202492428 & 488239798537030 & 39180748052700 & 1986889884620 & 5393481373052 \\ 0 & 6782207080 & 238370724380 & 2041791514140 & 8480781863740 & 208333346883670 & 332266439747106 & 34888681378940 & 24097462767010 & 1108868783760 & 28043348858044 \\ 57878637760 & 28564042298220 & 3098176511479280 & 14356440238205635 & 3991199974786200 & 6987865311080820 & 81762211401327198 & 6521352883777995 & 338328730973790 & 115312720216130 & 28509977559272 \\ 498282579760 & 252414547688728 & 194679747235820 & 7809258025046420 & 165724808432320 & 233328219344424 & 25297282028181852 & 1748441876085540 & 7870665648007900 & 21129116080908 & 377409107237628 \\ 82687442447120 & 146638748071388 & 896629728204120 & 2704029272288830 & 52500679384868280 & 67526748901207324 & 8485864724419364 & 33707337050110470 & 1263780830762420 & 2517631025243518 & 100310771884846 \\ 148140314822640 & 5881144622708016 & 27906262771514000 & 7240085337919890 & 198318687198670400 & 131071181160122648 & 9600891421160308 & 4523389406272600 & 1254982652811640 & 108242769642964 & 0 \\ 1754184701373740 & 1607733127380292 & 6335138217720660 & 1404602584860210 & 1972501870281840 & 1819808232340796 & 10380671394478244 & 386368920976090 & 501172233370481 & 0 & 0 \\ 439673140007600 & 31771283186568156 & 1035792171132940 & 193938078844401000 & 22662178817967640 & 1716811248944619828 & 768783125513892720 & 14726562829687989 & 0 & 0 & 0 \\ 70706027359440 & 42142727981133820 & 11662866391620320 & 187004174964081860 & 17633830408979200 & 10063921329237760 & 26672860281531376 & 0 & 0 & 0 & 0 \\ 712526958986440 & 3786625201053360 & 8577182988763280 & 1127420403129960 & 820524248199040 & 2776642514174840 & 0 & 0 & 0 & 0 & 0 \\ 4015182786100800 & 190178878072832800 & 371893024262944800 & 3747643864777600 & 1763962690692400 & 0 & 0 & 0 & 0 & 0 & 0 \\ 95728301094400 & 452632838792000 & 7271669102544000 & 496172949856400 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

#### 4. CONCLUDING REMARKS

1. In our computer algebra approach for derivation of bounds for the best Markov constant  $c_n(\alpha)$  we perform some optimization with respect to parameter  $s$ .

Our motivation for searching lower bounds for  $c_n^2(\alpha)$  with a factor depending on  $n$  of the special form  $n(n + \sigma(\alpha + 1))$  is Corollary D(ii).

An interesting observation about the lower bounds  $\underline{c}_{n,k}(\alpha)$  in Theorem 1 is that they imply

$$\frac{kn}{k+1} = \lim_{\alpha \rightarrow \infty} \alpha \underline{c}_{n,k}^2(\alpha) \leq \lim_{\alpha \rightarrow \infty} \alpha c_n^2(\alpha), \quad 3 \leq k \leq 6$$

(the lower bound in Corollary 2 follows from the case  $k = 6$ ). This observation and Proposition 3 give rise for the following

**Conjecture 1.** The best Markov constant  $c_n(\alpha)$  satisfies:

$$\lim_{\alpha \rightarrow \infty} \alpha c_n^2(\alpha) = n.$$

We also performed a search for lower bounds for  $c_n^2(\alpha)$  with a factor depending on  $n$  of the form  $(n+1)(n+\sigma(\alpha+1))$ . Such a choice is reasonable, as the resulting lower bounds preserve the limit relation in Corollary D (i). The optimal value then is  $\sigma = -1/3$  (the same for all  $k$ ,  $3 \leq k \leq 6$ ), and we obtain lower bounds as in Theorem 1 with  $n(n + \sigma(\alpha + 1))$  replaced by  $(n + 1)(n - (\alpha + 1)/3)$ . These lower bounds make sense only for  $n > (\alpha + 1)/3$ , and are better than those in Theorem 1 only for  $\alpha$  close to  $-1$ .

**2.** The bounds  $(\underline{c}_{n,k}(\alpha), \bar{c}_{n,k}(\alpha))$  ( $3 \leq k \leq 6$ ) in Theorem 1 imply bounds  $(\ell_k(\alpha), u_k(\alpha))$  (occurring in the middle columns of Tables 1 and 2) for the asymptotic Markov constant  $c(\alpha)$ , and the bounds deduced with a larger  $k$  are superior. While the lower bounds  $\ell_k(\alpha)$  are of the correct order  $\mathcal{O}(\alpha^{-1})$  as  $\alpha \rightarrow \infty$ , for the upper bound  $u_k(\alpha)$  we have  $u_k(\alpha) = \mathcal{O}(\alpha^{-1+\frac{1}{2k}})$  as  $\alpha \rightarrow \infty$ , ( $3 \leq k \leq 6$ ). The ratio

$$\rho_k(\alpha) := \frac{u_k(\alpha)}{\ell_k(\alpha)}, \quad 3 \leq k \leq 6,$$

tends to 1 as  $\alpha \rightarrow -1$ , which indicates that for moderate  $\alpha$  the bounds  $\ell_k(\alpha)$  and  $u_k(\alpha)$  are rather tight. This observation is clearly seen in the particular case  $\alpha = 0$ , where, according to Turán's result, we have  $c(0) = \frac{2}{\pi}$ . We give the lower and the upper bounds for  $c(0)$  and the overestimation factors in Table 3.

**3.** Another interesting observation, concerning the coefficients of  $R_n$  inspires the following

**Conjecture 2.** For every fixed  $k \in \mathbb{N}$ , the coefficient  $b_{k,n}$ ,  $n > k$ , of the polynomial  $R_n(x) = x^n - b_{1,n}x^{n-1} + b_{2,n}x^{n-2} - \dots + (-1)^n b_{n,n}$ , satisfies

$$b_{k,n} = \frac{n^{2k}}{2^k k!(\alpha+1) \cdots (\alpha+2k-1)} + \mathcal{O}(n^{2k-1}). \quad (4.1)$$

Conjecture 2 is verified with our computer algebra approach for  $1 \leq k \leq 6$ , but so far we do not have a proof for the general case. Having (4.1) proved,

we could try to find the explicit form of  $d_k$ , the coefficient of  $n^{2k}$  in Newton's function  $p_k(R_n)$ , and consequently to obtain two sequences  $\{\ell_k\}$  and  $\{u_k\}$  defined by  $\ell_k = \sqrt{d_k/d_{k-1}}$  and  $u_k = \sqrt[2^k]{d_k}$  which converge monotonically from below and from above, respectively, to  $c(\alpha)$ , the sharp asymptotic Markov constant.

Table 3: The lower and the upper bounds for the asymptotic Markov constant  $c(0)$  and the overestimation factors.

$k$	$\ell_k(0)$	$u_k(0)$	$\frac{c(0)}{\ell_k(0)}$	$\frac{u_k(0)}{c(0)}$
3	$\sqrt{\frac{2}{5}} \approx 0.63245553$	$\sqrt[6]{\frac{1}{15}} \approx 0.63677321$	1.006584242	1.00024103
4	$\sqrt{\frac{17}{42}} \approx 0.63620901$	$\sqrt[8]{\frac{17}{630}} \approx 0.63663212$	1.00064564	1.00001939
5	$\sqrt{\frac{62}{153}} \approx 0.63657580$	$\sqrt[10]{\frac{31}{2835}} \approx 0.63662085$	1.00006906	1.00000170
6	$\sqrt{\frac{2073}{5115}} \approx 0.63661494$	$\sqrt[12]{\frac{2073}{467775}} \approx 0.63661987$	1.00000757	1.00000015

Although the ratios  $\rho_k$ ,  $3 \leq k \leq 6$ , satisfy  $\rho_k(\alpha) \rightarrow \infty$  as  $\alpha \rightarrow \infty$ , they grow rather slowly. For instance,  $\rho_6(\alpha) < 2$  for  $\alpha < 140000$ , see Figure 1.

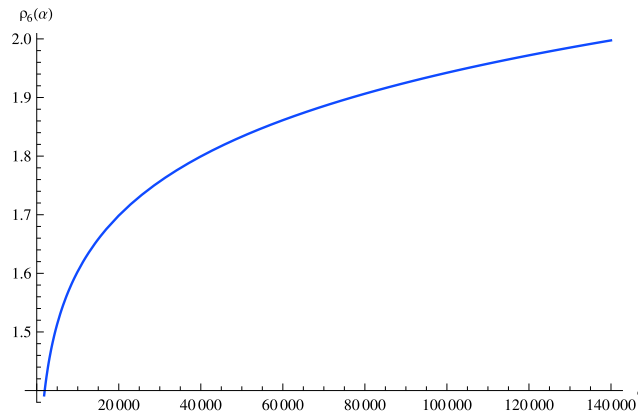


Figure 1: The graph of  $\rho_6(\alpha) < 2$ .

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